Multistage Games with Payoff Uncertainty

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Abstract

This lecture introduces the theory of multistage games with observable actions and either incomplete information, or asymmetric information about an initial chance move. It is focused on the description of these games and on the role of Bayes rule to model updated beliefs about co-players' types. The analytical representation blends the parameterized payoff functions (with differential information about parameters) used to study static games with incomplete information in Chapter 8 with the sequential game-tree representation of multistage games of Chapter 9. Graphically, the game is represented as a forest of trees, one for each possible profile of payoff functions, connected by information sets showing that (active) players cannot distinguish the different types of co-players. Conditional probability systems about the types and strategies of co-players yield "personal assessments" comprising conditional probabilities of actions and conditional probabilities of nodes in the information sets. Personal assessments must satisfy Bayes rule.

[These slides summarize and, in part, complement Sections 1, 2 and part of section 3 of Chapter 15 of GT-AST.]

Introduction

- We want to study multistage games with observable actions and incomplete information.
 - The same analytical tools can be applied to study games with complete, but imperfect and asymmetric information, e.g., Poker.
- To represent incompleteness of information in static games:
 - we used parameterized payoff functions $(u_i : \Theta \times A \rightarrow \mathbb{R})_{i \in I}$,
 - $\theta = (\theta_i)_{i \in I} \in \Theta$, where each $i \in I$ knows (only) θ_i (for simplicity, we neglect residual uncertainty).
- To represent the rules of interaction of a **multistage game with observable actions**:
 - we used a structure ⟨I, (A_i, A_i (·))_{i∈I}⟩ with A_i=i's potentially feasible actions, A_i (·)=i's constraint correspondence;
 - we derived a **tree** of histories (\bar{H}, \preceq) , with $\preceq =$ "prefix of", and $\bar{H} = H \cup Z$ (Z: terminal histories).
- To represent *incomplete information* (payoff uncertainty), posit (u_i : Θ × Z → ℝ)_{i∈I}.

Graphical representation: a simple example

- We get a *forest of game trees* indexed by θ ∈ Θ, each one with payoff functions (u_{i,θ} : Z → ℝ)_{i∈I}.
- Asymmetric information is graphically represented by information sets: for every *i* ∈ *I*, *θ*_i ∈ Θ_i, and *h* ∈ *H*, the *indistinguishable* "nodes" ((*θ*_i, *θ*_{-i}), *h*) (*θ*_{-i} ∈ Θ_{-i}) are connected by a dashed line. In Figure A, *i* = 2, Θ₂ = {*θ*₂}, Θ₁ ≅ Θ = {*θ*', *θ*''}, *h* = (*r*).



Intuitive analysis of the example

In Figure A we also represent the initial, exogenous belief of pl. 2 about θ. His (posterior) belief upon observing r may well be different: r is dominated for θ", but it is justifiable for θ'; if 2 strongly believes in 1's rationality, then the posterior belief is μ₂ (θ"|r) = 1 − μ₂ (θ'|r) = 0.



- Also, d is conditionally dominated by u. A rational pl. 2 would choose u after r. Anticipating this, a rational player 1 of type θ' chooses r.
- Intuitive solution: $(r'.\ell'', u)$ [where $r'.\ell'' = (r \text{ if } \theta'', \ell \text{ if } \theta'')$].

Multistage games with payoff uncertainty

A **multistage game with payoff uncertainty** (and observable actions) is a structure

$$\hat{\Gamma} = \left\langle I, (\Theta_i, A_i, \mathcal{A}_i(\cdot), u_i)_{i \in I} \right\rangle$$

where

- $\langle I, (A_i, \mathcal{A}_i(\cdot))_{i \in I} \rangle$ is a multistage game tree;
- $\theta_i \in \Theta_i$ are information **types** (traits, tastes, private information);
- $\Theta := \times_{i \in I} \Theta_i$ (more generally, $\Theta \subseteq \times_{i \in I} \Theta_i$);
- $u_i: \Theta \times Z \to \mathbb{R}$ is *i*'s parameterized **payoff function**;
 - u_i is derivable from $g: \Theta \times Z \to Y$ and $v_i: \Theta_i \times Y \to \mathbb{R}$;
 - generalization: A_i(·, ·) : Θ_i × H ⇒ A_i, a_i ∈ A_i(θ_i, h) feasible for θ_i given h;
 - $\Theta = \{\overline{\theta}\}$ singleton (or θ -indep. payoff) \Rightarrow complete information.
- To ease notation, we neglect residual uncertainty (Θ₀ is a singleton, not shown).

An example



• $\Theta = \{\theta', \theta''\} \cong \Theta_1$, only pl. 1 is informed.

- Dashed lines join "nodes" (θ, h) in the same information set of the active player (pl. 2 and 3 have the same information).
- (Numbers in $\{\cdot\}$: exogenous common prior on Θ , not part of \hat{F})

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• **Mini-Poker:** θ_i is acquired information (not personal trait).

- $\Theta_i = \{J, Q, K\}, \Theta = \{JK, JQ, KQ, KJ, QJ, QK\} \subset \Theta_1 \times \Theta_2.$
- Initial node of pl. 0=chance move (random order of cards).

Personal Assessments

- For simplicity, focus on two-person games.
 - With complete information, we could represent a conjecture of *i* as $\beta^{i} = (\beta^{i}(\cdot|h))_{h\in H} = \beta_{-i} \in \times_{h\in H}\Delta(\mathcal{A}_{-i}(h))$, which is like a behavior strategy of -i, and we obtained a subjective decision tree for *i*.
 - With incomplete info, we need a personal assessment (βⁱ, μ_i):
 conjecture βⁱ must have the form

$$\beta^{i} = \left(\beta^{i}\left(\cdot|\theta_{-i},h\right)\right)_{\theta_{-i}\in\Theta_{-i},h\in H} \in \left(\times_{h\in H}\Delta\left(\mathcal{A}_{-i}\left(h\right)\right)\right)^{\Theta_{-i}},$$

because pl. *i* of type θ_i expects that -i's behavior depends of his type θ_{-i} ; also, we have to specify updated beliefs about θ_{-i} given each $h \in H$: a system of beliefs

$$\mu_{i} = (\mu_{i}(\cdot|h))_{h \in H} \in (\Delta(\Theta_{-i}))^{H}$$

(we also write μ_i (·|θ_i, h) to emphasize that θ_i is known/given).
Note: μ_i should be derived from βⁱ (if possible) using Bayes rule.
From (βⁱ, μ_i) we will derive conditional values for type θ_i.

Conditional Probability Systems (CPSs)

- In the (rationalizability) analysis of static games with incomplete information, we considered conjectures μⁱ ∈ Δ (Θ_{-i} × A_{-i}).
- In the (rationalizability) analysis of *multistage* games with complete information, we considered CPSs μⁱ ∈ Δ^H (S_{-i}).
- In the (rationalizability) analysis of *multistage* games with *incomplete* information, we can use CPSs $\bar{\mu}^i \in \Delta^H (\Theta_{-i} \times S_{-i})$, where—as before— $S_{-i} = \times_{h \in H} \mathcal{A}_{-i}(h)$ are the co-player's pure strategies (we write $\bar{\mu}^i$ to distinguish from systems of beliefs $\mu_i \in (\Delta(\Theta_{-i}))^H)$.
- We can derive a personal assessment (β^{i}, μ_{i}) from a CPS $\overline{\mu}^{i}$: $\forall (\theta_{-i}, h) \in \Theta_{-i} \times H, \forall a_{-i} \in \mathcal{A}_{-i}(h),$ $\mu_{i}(\theta_{-i}|h) = \overline{\mu}^{i}(\{\theta_{-i}\} \times S_{-i}(h)|h), \text{ if } \overline{\mu}^{i}(\{\theta_{-i}\} \times S_{-i}(h)|h) > 0$ then $\overline{\mu}^{i}(\{\theta_{-i}\} \times S_{-i}(h)|h), \text{ if } \overline{\mu}^{i}(\{\theta_{-i}\} \times S_{-i}(h)|h) > 0$

$$\beta^{i}\left(\mathbf{a}_{-i}|\theta_{-i},h\right) = \frac{\bar{\mu}^{\prime}\left(\left\{\theta_{-i}\right\} \times S_{-i}\left(h,a_{-i}\right)|h\right)}{\bar{\mu}^{i}\left(\left\{\theta_{-i}\right\} \times S_{-i}\left(h\right)|h\right)}$$

Example of Personal Assessment: Signaling



•
$$\mu_2(\theta'|\varnothing) = p(\theta') = 1/2,$$

• $\beta^2(r|\theta') > 0, \ \beta^2(r|\theta'') = 0,$
• $\mu_2(\theta'|r) = \frac{\beta^2(r|\theta')p(\theta')}{\beta^2(r|\theta')p(\theta')+\beta^2(r|\theta'')p(\theta'')} = 1.$

Example of Personal Assessment: Poker



•
$$p(\theta) = 1/6$$
 for each $\theta \in \{JK, JQ, KQ, KJ, QJ, QK\}.$
• $\beta^2(B|K_1) = \beta^2(B|Q_1) = 1, \beta^2(B|J_1) = q \in (0, 1).$
• $\mu_2(J_1|Q_2, B) = \frac{1}{6}q/(\frac{1}{6}q + \frac{1}{6}) = q/(q + 1).$

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Rehearsing Bayes Rule: Preliminaries

- For the first time, we are truly using Bayes rule!
- Abstract finite setting: uncertainty space

$$\Omega = \Theta \times X$$
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e.g., $\theta \in \Theta$ is the composition of a *k*-color *n*-balls urn, $x \in X$ is a sequence of colors of balls drawn from the urn.

- The likelihood function θ → P(·|θ) ∈ Δ(X) gives the probability law for x ∈ X given each parameter value θ ∈ Θ.
- As in *Bayesian statistics*, we posit $P(\cdot) \in \Delta(\Theta \times X)$ and interpret $P(x|\theta) = P(\{(\theta, x)\} | \{\theta\} \times X)$, where

$$P(\{(\theta, x)\} | \{\theta\} \times X) = \frac{P(\{(\theta, x)\})}{P(\{\theta\} \times X)} =: \frac{P(\theta, x)}{P(\theta)}$$

(recall: $P(E|F) = P(E \cap F|F) = P(E \cap F)/P(F)$ if P(F) > 0).

Rehearsing Bayes Theorem (or Rule)

- The joint probability of θ and x can be expressed as $P(x|\theta) P(\theta) = P(\theta, x) = P(\theta|x) P(x).$
- The marginal, or **predictive** probability of x is

$$P(x) := P(\Theta \times \{x\}) = \sum_{\theta' \in \Theta} P(\theta', x) = \sum_{\theta' \in \Theta} P(x|\theta') P(\theta').$$

• Therefore (**Bayes theorem**) the "posterior" probability of any $\theta \in \Theta$ given any evidence $x \in X$ is

$$P(\theta|x) = \frac{P(\theta, x)}{P(x)} = \frac{P(x|\theta) P(\theta)}{\sum_{\theta' \in \Theta} P(x|\theta') P(\theta')}.$$

- (1) Bayes rule only uses the likelihood function $\theta \mapsto P(\cdot|\theta)$ and the "prior" $\operatorname{marg}_{\Theta} P \in \Delta(\Theta)$.
- (2) In statistics, we may assume P(x) > 0, not in GT where x is endogenous.

Bayes Rule: Connection to Games

- Take the perspective of an external observer of a multistage game with incomplete information.
- The "probability law" is $\beta = (\beta (\cdot | \theta, h))_{\theta \in \Theta, h \in H} \in (\times_{h \in H} \Delta (\mathcal{A} (h)))^{\Theta} \text{ (obtained from a profile of type-dependent behavior strategies).}$
- There is an exogenous prior $p \in \Delta(\Theta)$; write $\beta(a|\theta, \emptyset) = \beta(a|\theta)$; for each $a \in \mathcal{A}(\emptyset)$ such that $\sum_{\theta' \in \Theta} \beta(a|\theta') p(\theta') > 0$ and each $\theta \in \Theta$,

$$\mu\left(\theta|\mathbf{a}\right) = \frac{\beta\left(\mathbf{a}|\theta\right)p\left(\theta\right)}{\sum_{\theta'\in\Theta}\beta\left(\mathbf{a}|\theta'\right)p\left(\theta'\right)}.$$

• Let $\mu(\cdot|h) \in \Delta(\Theta)$ be given; for each $a \in \mathcal{A}(h)$ such that $\sum_{\theta' \in \Theta} \beta(a|\theta', h) \mu(\theta|h, a) > 0$ and each $\theta \in \Theta$,

$$\mu(\theta|h, \mathbf{a}) = \frac{\beta(\mathbf{a}|\theta, h) \mu(\theta|h)}{\sum_{\theta' \in \Theta} \beta(\mathbf{a}|\theta', h) \mu(\theta'|h)}$$

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