

LECTURES 8

Solving models with occasionally binding constraints

Part II

Macroeconomics 4

A.Y. 2014-15

Time iteration

- Consider the Euler equation for our recursive problem:

$$u_c(c_{z,i}) - \varphi_{z,i} = \beta(1+r) \sum_{j=1}^n \Pi_{ij} u_c(c'_{z,i,j}),$$

$$c_{z,i} = (1+r)a_z + ws_i - a'_i(a_z),$$

$$c'_{z,i,j} = (1+r)a'_i(a_z) + ws_j - a'_j[a'_i(a_z)],$$

$$a'_i \geq 0,$$

$$\varphi_{z,i} a'_i = 0,$$

$$\varphi_{z,i} \geq 0.$$

where $z \in \{1, 2, \dots, m\}$ and $i, j \in \{1, 2, \dots, n\}$.

Time iteration

- Suppose that a current guess for the policy function(s), say $a'_{k,i}(a)$, exists.
- For the moment, set $\varphi_{z,i} = 0$ for all z and i . The Euler eq. boils down to:

$$u_c [(1+r)a_z + ws_i - h_{k,i}(a_z)] = \beta(1+r) \times \sum_{j=1}^n \Pi_{ij} u_c \left\{ (1+r)h_{k,i}(a_z) + ws_j - a'_{k,j}[h_{k,i}(a_z)] \right\}.$$

- This is a system of $m \times n$ non-linear eqs. that can be solved for the values of $h_{k,i}(a_z)$, using some robust method like the *Trust Region algorithm*.
- Note that the term $a'_{k,j}[h_{k,i}(a_z)]$ has to be computed via interpolation.

Time iteration

- The current guess $a'_{k,i}(a)$ is then updated in the following way:

$$\begin{cases} a'_{k+1,i}(a_z) = 0 & \text{if } h_{k,i}(a_z) \leq 0 \\ a'_{k+1,i}(a_z) = h_{k,i}(a_z) & \text{if } h_{k,i}(a_z) > 0 \end{cases}, \quad \forall i, z.$$

- Convergence can be assessed with the usual stopping rule:

$$\|a'_{k+1,i}(a) - a'_{k,i}(a)\| \leq \varepsilon > 0, \quad \forall i.$$

- Note that, if needed, the multiplier $\varphi_{z,i}$ can be obtained from:

$$\varphi_{z,i} = u_c(c_{z,i}) - \beta(1+r) \sum_{j=1}^n \Pi_{ij} u_c(c'_{z,i,j}).$$

Time iteration

- This method is known as *Time Iteration on the Euler equation* (TI).
- In standard models, i.e. without borrowing constraints, under the usual regularity conditions the method converges to a solution of the Bellman equation.
- Rendhal (2013) shows that is the case also for models with occasionally binding constraints.
- He deals with problems that can be formulated as:

$$V(a, s) = \max_{a' \in \Gamma(a, s)} \{F(a, a', s) + \beta \mathbb{E}[V(a', s') \mid s]\},$$

where $F(\cdot)$ is **bounded** over a **compact** set A , and:

$$\Gamma(a, s) = \{a' \in A : m_j(a, a', s) \leq 0, j = 1, 2, \dots, r\}.$$

Time iteration

- Standard results show that, under some regularity conditions, the *VFI* scheme:

$$V_{n+1}(a, s) = \max_{a' \in \Gamma(a, s)} \{F(a, a', s) + \beta \mathbb{E}[V_n(a', s') | s]\},$$

$$a'_{n+1}(a, s) = \arg \max_{a' \in \Gamma(a, s)} \{F(a, a', s) + \beta \mathbb{E}[V_n(a', s') | s]\},$$

will uniformly converge, respectively, to the unique fixed points V and a' for any weakly concave and bounded V_0 .

Time iteration

- The previous *VFI* scheme can be written as:

$$V_{n+1}(a, s) = \min_{\boldsymbol{\mu} \geq 0} \max_{a' \in A} L(a, a', s, \boldsymbol{\mu}) =$$

$$\max_{a' \in A} \min_{\boldsymbol{\mu} \geq 0} L(a, a', s, \boldsymbol{\mu}),$$

where:

$$L(a, a', s, \boldsymbol{\mu}) = F(a, a', s) + \beta \mathbb{E}[V_n(a', s') | s] - \boldsymbol{\mu}^T \mathbf{m}(a, a', s),$$

and $\boldsymbol{\mu}$ is a $r \times 1$ column vector of Lagrange multipliers.

Time iteration

- The *TI* scheme can be described as:

$$F_{a'}(a, a', s) - \mathbf{J}_{a'}(a, a', s) \boldsymbol{\mu}_{k+1}(a, s) + \\ \beta \mathbb{E} [F_a[a', h_k(a', s'), s'] | s] + \\ - \beta \mathbb{E} [\mathbf{J}_a(a', h_k(a', s'), s') \boldsymbol{\mu}_k(a', s') | s] = 0$$

where $a' = h_{k+1}(a, s)$, and \mathbf{J}_x the Jacobian of \mathbf{m} w.r.t. x .

- Rendhal (2013) proves that, if $h_k(a, s) = a'_k(a, s)$, then:

$$h_{k+1}(a, s) = a'_{k+1}(a, s).$$

- Hence, $h_k(a, s) \rightarrow a'(a, s)$ when $k \rightarrow \infty$: *TI* converges to the true solution also under occasionally binding constraints.

Fixed point iteration

- *TI* has solid theoretical foundations, but is still cumbersome to implement, and relatively slow to converge.
- A simple variant, known as *Fixed Point Iteration (FPI)*, has some useful properties, and can be characterized as:

$$u_c [(1+r) a_z + w s_i - h_{k,i}(a_z)] = \beta (1+r) \times \sum_{j=1}^n \Pi_{ij} u_c \left\{ (1+r) a'_{k,i}(a_z) + w s_j - a'_{k,j} [a'_{k,i}(a_z)] \right\}.$$

- Note that the r.h.s. is computed using the current guess $a'_{k,i}(a_z)$ only, so there is no need to solve any non-linear system: you just need to invert u_c to solve for $h_{k,i}(a_z)$.
- The guess is updated with the same rule as before:

$$\bar{a}'_{k,i}(a_z) = \max [h_{k,i}(a_z), 0].$$

Fixed point iteration

- There aren't theoretical results on the convergence properties of *FPI*, to the best of my knowledge.
- However, convergence can be improved, or even “forced,” by using a simple **damping scheme**.
- In other words, the updating rule becomes:

$$\begin{cases} a'_{k+1,i}(\mathbf{a}_z) = 0 & \text{if } h_{k,i}(\mathbf{a}_z) \leq 0 \\ a'_{k+1,i}(\mathbf{a}_z) = \varrho \bar{a}'_{k,i}(\mathbf{a}_z) + (1 - \varrho) a'_{k,i}(\mathbf{a}_z) & \text{if } h_{k,i}(\mathbf{a}_z) > 0 \end{cases}'$$

for some $\varrho \in (0, 1)$, say $\varrho = 0.9$.

- There are no general rules - some nasty cases require low values of ϱ to converge - but *FPI* is considerably faster, and, if it converges, then it converges to the unique solution of the *VFI* scheme.

Fixed point iteration

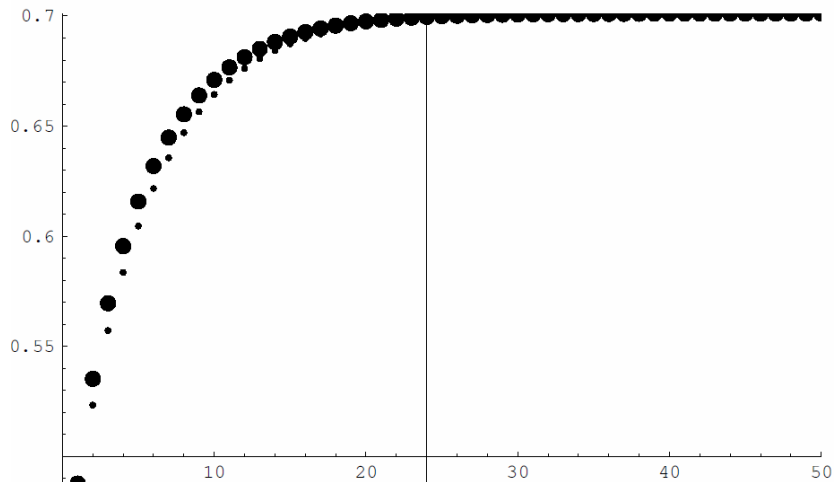


Figure 4 $c_{T-s}(k, \theta)$ Using Time Iteration versus Fixed Point Iteration

References I

Rendhal, P. (2013). Inequality Constraints and Euler Equation based Solution Methods. *The Economic Journal*, forthcoming.